

LogChow Theory.

Use log language throughout

Motivation: Given $X \xrightarrow{f} Y$ a morphism we can construct

functorial Grothendieck maps $f^!: A_*(Y) \rightarrow A_*(X)$ by

embedding the intrinsic normal cone \mathbb{C}_f into E a

vector bundle stack. This gives refined maps

for any diagram

$$\begin{array}{ccc} W & \xrightarrow{f_1} & V \\ \downarrow & \lrcorner & \downarrow \\ X & \xrightarrow{f} & Y \end{array} \quad f_{E,V}^!: A_*(V) \rightarrow A_*(W).$$

i.e. an invariant class $f_E^! \in A_*(X \xrightarrow{f} Y)$.

If instead $f: X \rightarrow Y$ is a log morphism all this
breaks, log fibre products don't agree with
scheme fibre products etc.

One way to make this work (due with restrictions
by Leo Herr) would be to pull classes back from
Log $^+$, Olszak's log stack using the log normal
cone $\mathbb{C}_{xy}^+ = \mathbb{C}_{xy_{\log}}^+$.

Here's a way to generalise this:

$$X \xrightarrow{\text{Log}} \text{Log} \rightarrow Y$$

↑
log et

Thm (Fulton-Macpherson): Let $g: Y \rightarrow Z$ be smooth
relative dimension d and $f: X \rightarrow Y$ a map then

$$\bullet [g]: A^*(X \xrightarrow{f} Y) \xrightarrow{\sim} A^{*-d}(X \xrightarrow{f \circ g} Z)$$

So if the above theorem held for log smooth
morphisms we could construct using the strict
map $X \rightarrow \text{Log}$, and "pullback" to $X \rightarrow Y$.

The theorem doesn't hold for log smooth things.

But it hints that we at least need to produce
groups invariant under log blowups, (i.e. use
smoothing like b-Chow).

§ : Log Geometry

We've seen some of these definitions before in other
talks, so I'll try to be brief.

A log structure on X is:

morphism of sheaves of monoids

$M_X \xrightarrow{\alpha_X} \mathcal{O}_X^*$

sheaf of monoids
in et topology

multiplicative
monoid str.

such that $\alpha_X: \mathcal{O}_X^* \rightarrow \mathcal{O}_X^*$ is an isomorphism.

We have $\bar{M}_X = M_X / \mathcal{O}_X^*$.

fs \Rightarrow étale locally we have:

$$U \xrightarrow{\text{dru}} \text{Spec } k[\bar{P}], M_P \leftarrow \text{toric log str}$$

ét ↓

X

such that dru identifies $(M_P, \mathcal{O}_{\text{Spec } k[\bar{P}]})_u$ with

$$\text{dru}(M_P, \mathcal{O}_{\text{Spec } k[\bar{P}]})_u$$

K & Pa universal

ideal

eg $P^1 \xrightarrow{\text{dru}} \mathbb{A}^2$

$$\mathbb{A}^1 \xrightarrow{\text{dru}} \mathbb{A}^2$$

P^1 has non-toric

log str.

Here π_x has a degree zero blowup. In general

if $\mathcal{M}_{x,x}$ and $\mathcal{M}_{\tilde{x},\tilde{x}}$ are isomorphic to \mathbb{N}^k and $\mathbb{N}^{k'}$

for all x and \tilde{x} there are local blowups giving to a global degree zero map (have to apply

results of Cheung-Liu to lift from ét top). We say

that X and \tilde{X} (or instead just π_x) is locally free.

Dfn: $A_*(X) = \varinjlim A_*(\tilde{X})$ over all $\tilde{X} \rightarrow X$ a log blowup with \tilde{X} locally free.

§ Push and Pull.

Example that care is needed:

$$\begin{array}{ccc} P' & \xrightarrow{\pi'} & P \\ \pi' \downarrow & & \downarrow \pi' \\ P' & \longrightarrow & N^2 \\ & \pi' & \end{array}$$

does not commute for dimension reasons. The key things needed are having compatible fundamental class and integrality.

Dfn: Compatible fundamental class: X has compatible fundamental class if X is locally free and for all locally free maps

$\tilde{X} \xrightarrow{\pi'_x} X$ there is an equality $\pi'_x! [\tilde{X}] = [\tilde{X}]$

- * Lemma not stated in the talk: For any locally free X
- * there is a log blowup $\pi_x: \tilde{X} \rightarrow X$ and $V \subset \tilde{X}$ strict
- * such that: $[V] = \pi'_x! [X]$ and V has c.f.c. "wiring lemma"

For definition of integral see e.g. Neilt/Ogus.

Thm: (F.Kato) if $f: X \rightarrow Y$ is a log isomorphism then there is a diagram:

$$\begin{array}{ccc} \tilde{X} & \xrightarrow{f} & \tilde{Y} \\ \downarrow \pi_x & & \downarrow \pi_y \\ X & \xrightarrow{f} & Y \end{array} \quad \begin{array}{l} \pi_y \text{ a log blowup} \\ \pi_x \text{ a log modification} \\ f \text{ integral} \end{array}$$

So we can define push and pullbacks by looking at diagrams.

$$\begin{array}{ccc} \tilde{X} & \xrightarrow{f} & Y \\ \downarrow & & \downarrow \\ X & \xrightarrow{f} & Y \end{array} \quad \begin{array}{l} \bullet \text{ integral} \\ \bullet \text{ log blowups} \end{array}$$

and use blowups to move around. You get the usual compatibilities except only saturated pushforward and log flat pullback commute. Nonetheless this is enough to prove the original theorem.

Key compatibility:

Thm 3.2: Given $f: X \xrightarrow{f} Y$ such that Y has c.f.c. and

$\pi: \tilde{X} \rightarrow X$ a locally free log blowup then

$$i: \mathcal{C}_{\tilde{X}/X}^l \hookrightarrow \mathcal{C}_{X/Y}^l \times_X \tilde{X} \text{ and } i_*([\mathcal{C}_{\tilde{X}/X}^l]) = \pi^! ([\mathcal{C}_{X/Y}^l])$$

note that this does not happen naively for dimension reasons:

$$P' \rightarrow \mathrm{Bl}_A'$$

$$\downarrow \quad \downarrow$$

$$N^2 \cdot \hookrightarrow A^2$$